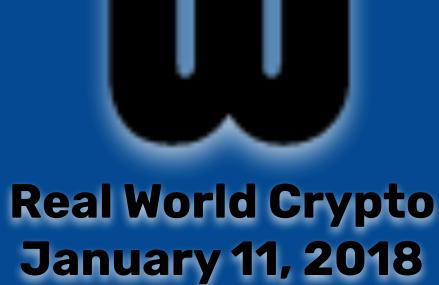


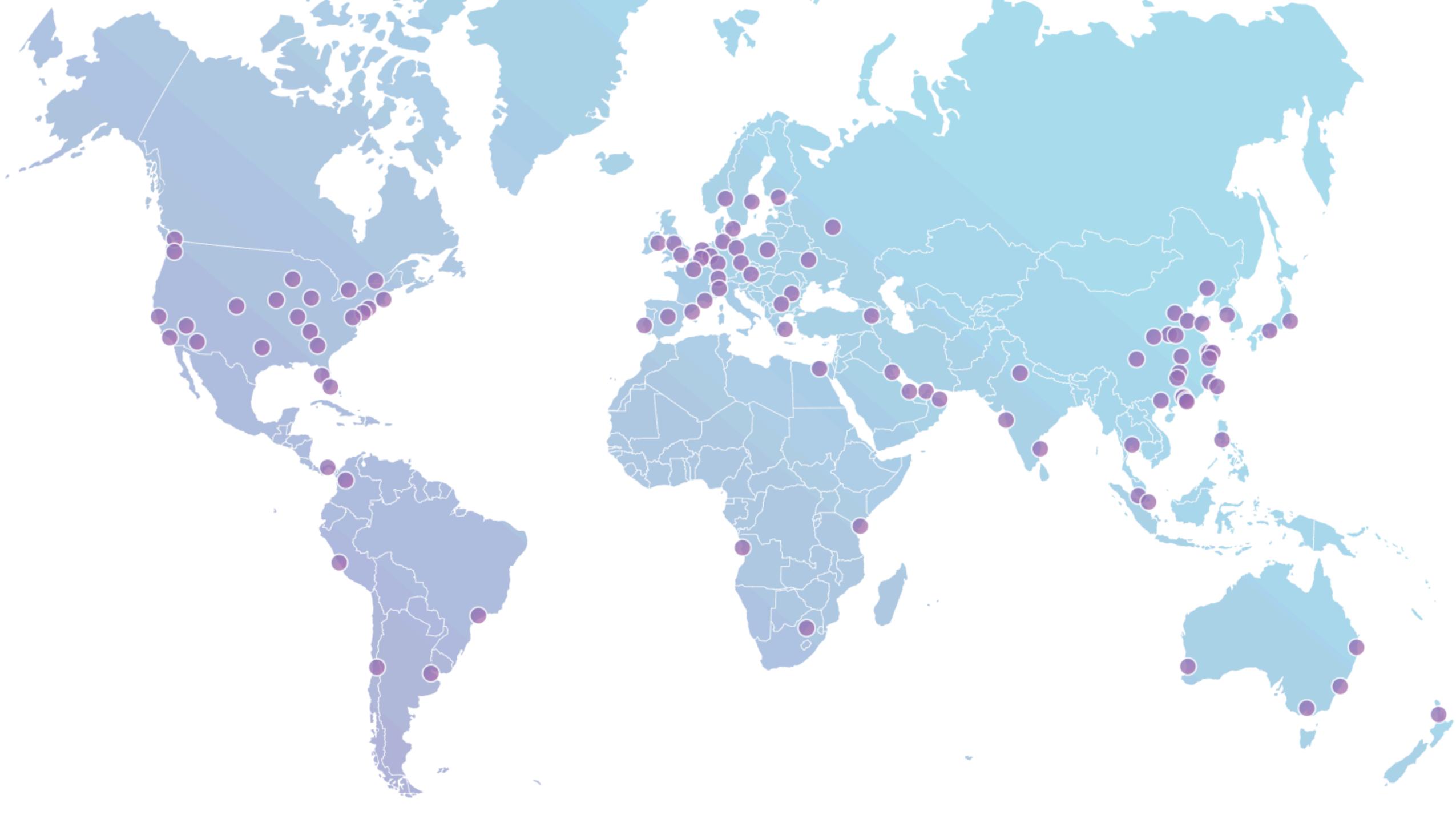
#### **Geo Key Manager Nick Sullivan** (@grittygrease) Brendan McMillion





# Geographically-Distributed Key Management



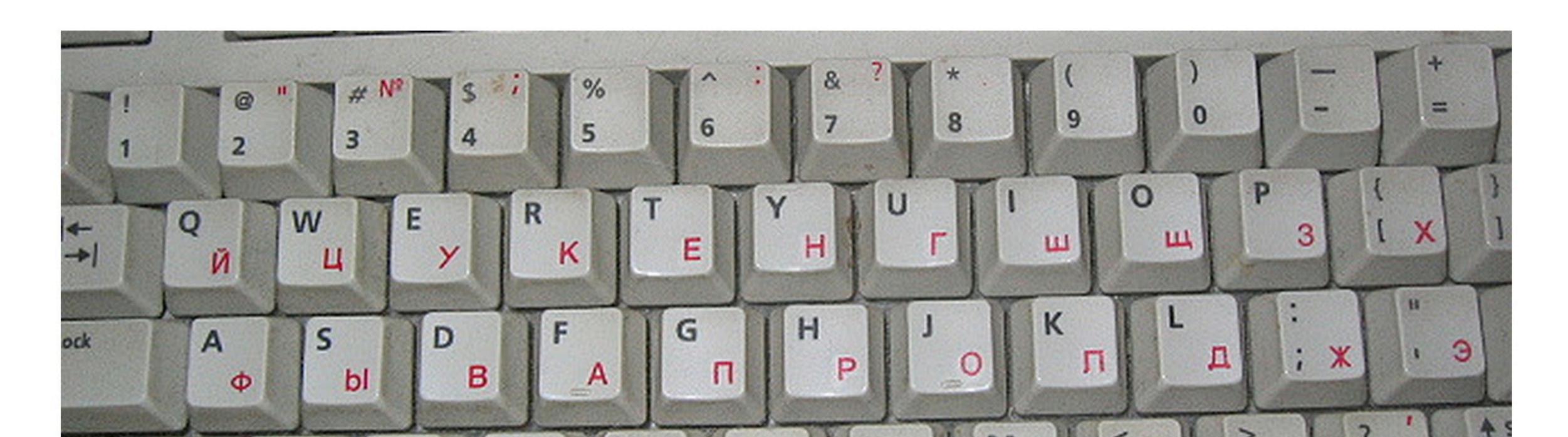




Layer 8

#### **Russia begins collecting encryption** keys while internet companies, like Facebook, stay silent

Patrick Howell O'Neill—July 27, 2016 at 4:27PM | Last updated July 27, 2016 at 9:29PM



#### The Daily Dot



So far, WhatsApp, Viber, and Telegram haven't said a public word.



Some things we have never done

- Cloudflare has never turned over our SSL keys or our customers' SSL keys to anyone.
- Cloudflare has never installed any law enforcement software or equipment anywhere on our network. •
- Cloudflare has never terminated a customer or taken down content due to political pressure.
- Cloudflare has never provided any law enforcement organization a feed of our customers' content transiting our network.

If Cloudflare were asked to do any of the above, we would exhaust all legal remedies, in order to protect our customers from what we believe are illegal or unconstitutional requests.



# Customer's choice Choose where in the world their keys are kept

# Deployability Work within existing constraints

Support network expansion





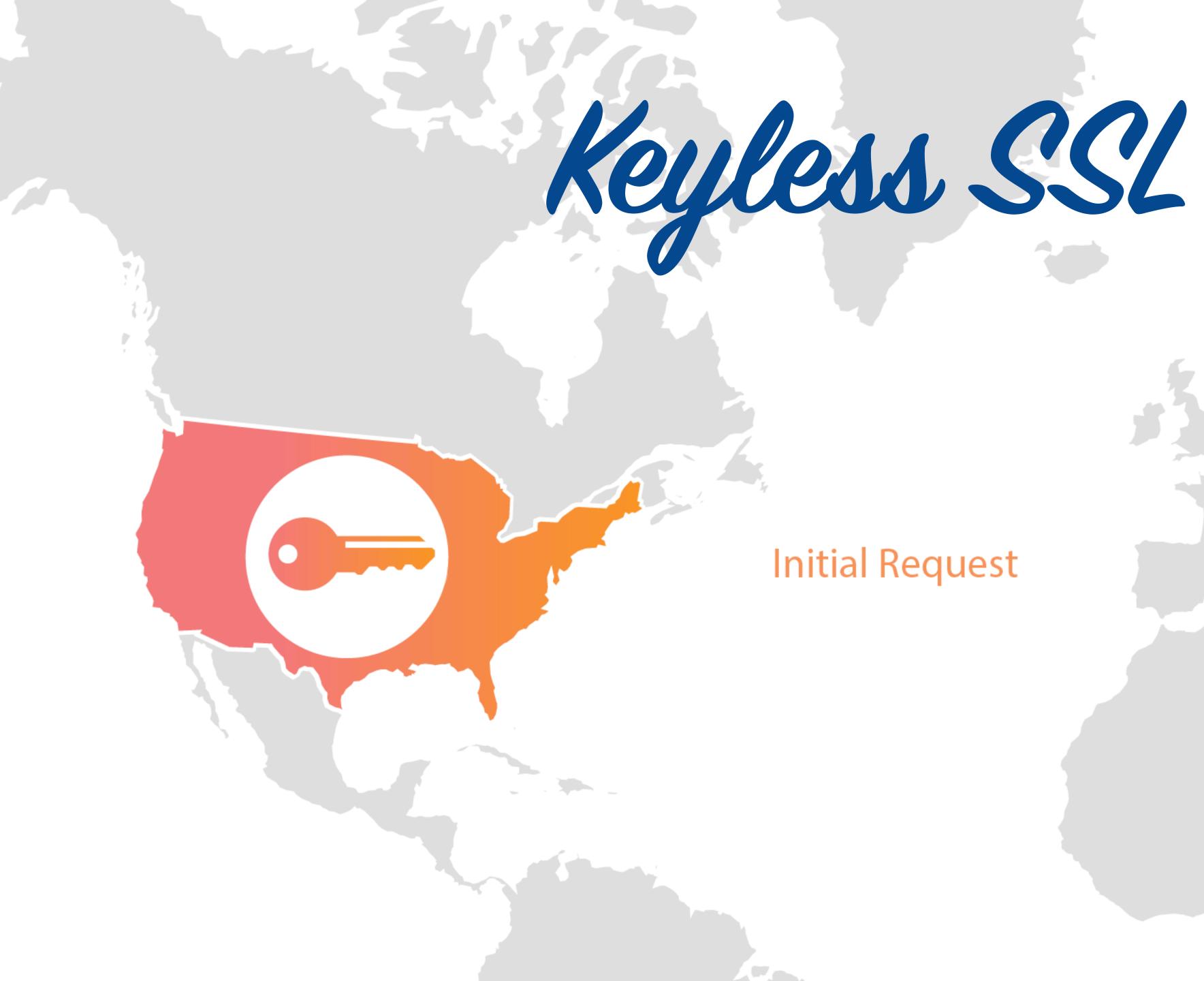
#### Constraint

Legacy client software



#### Component

### Keyless SSL







Amsterdam to Dusseldorf Los Angeles to Belgrade





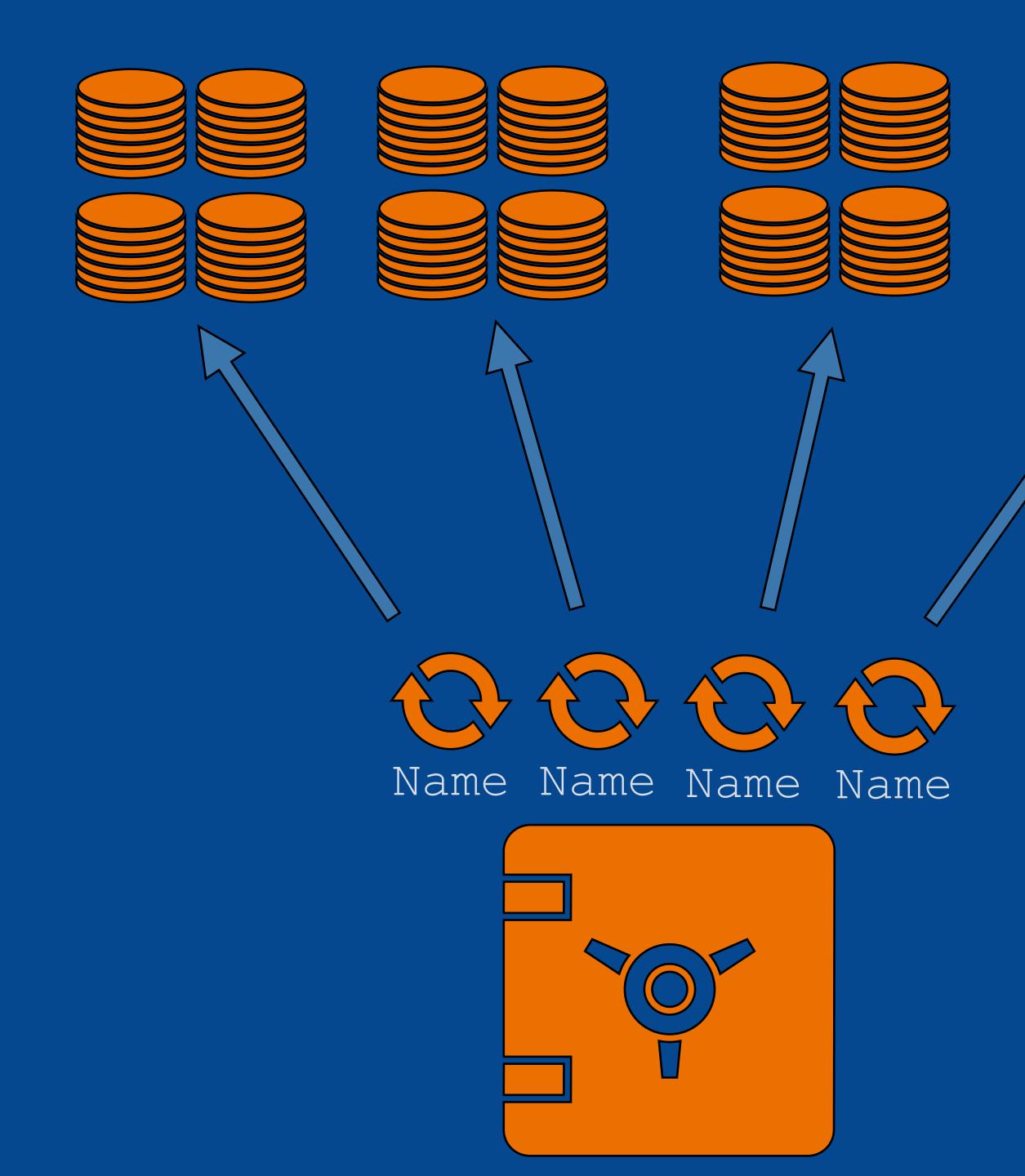
### Provisioning

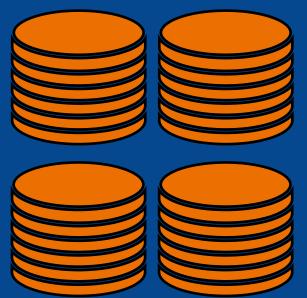






11





### Edge Machines

#### Template

#### Provisioning Server



#### Component

#### Provisioning

System



Constraints

# Non-interactive, Identity-based

#### Component

# Globally Synchronized Database





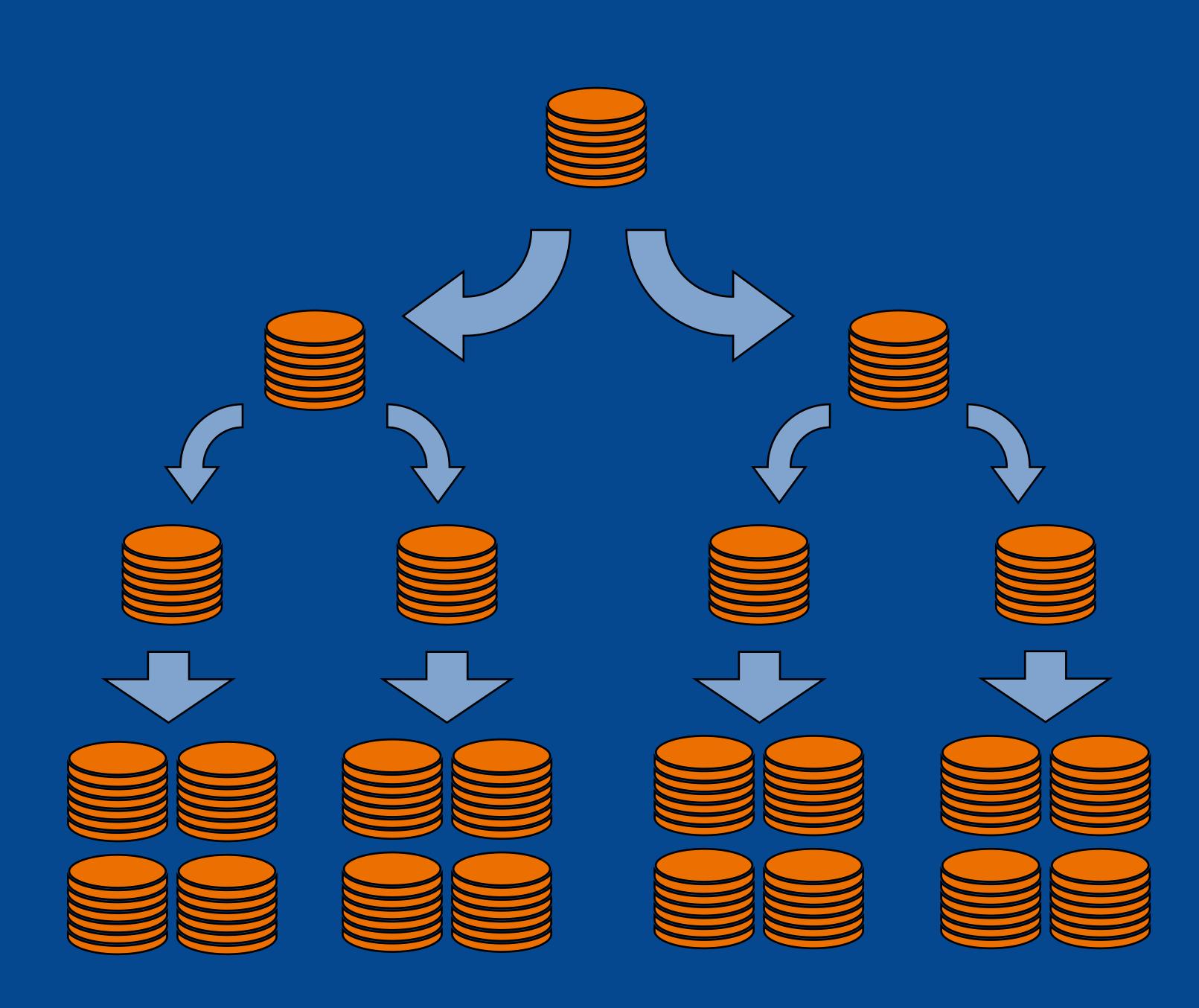


#### Master Database

#### Regional Master

#### Location Master

Local Copy



#### Component

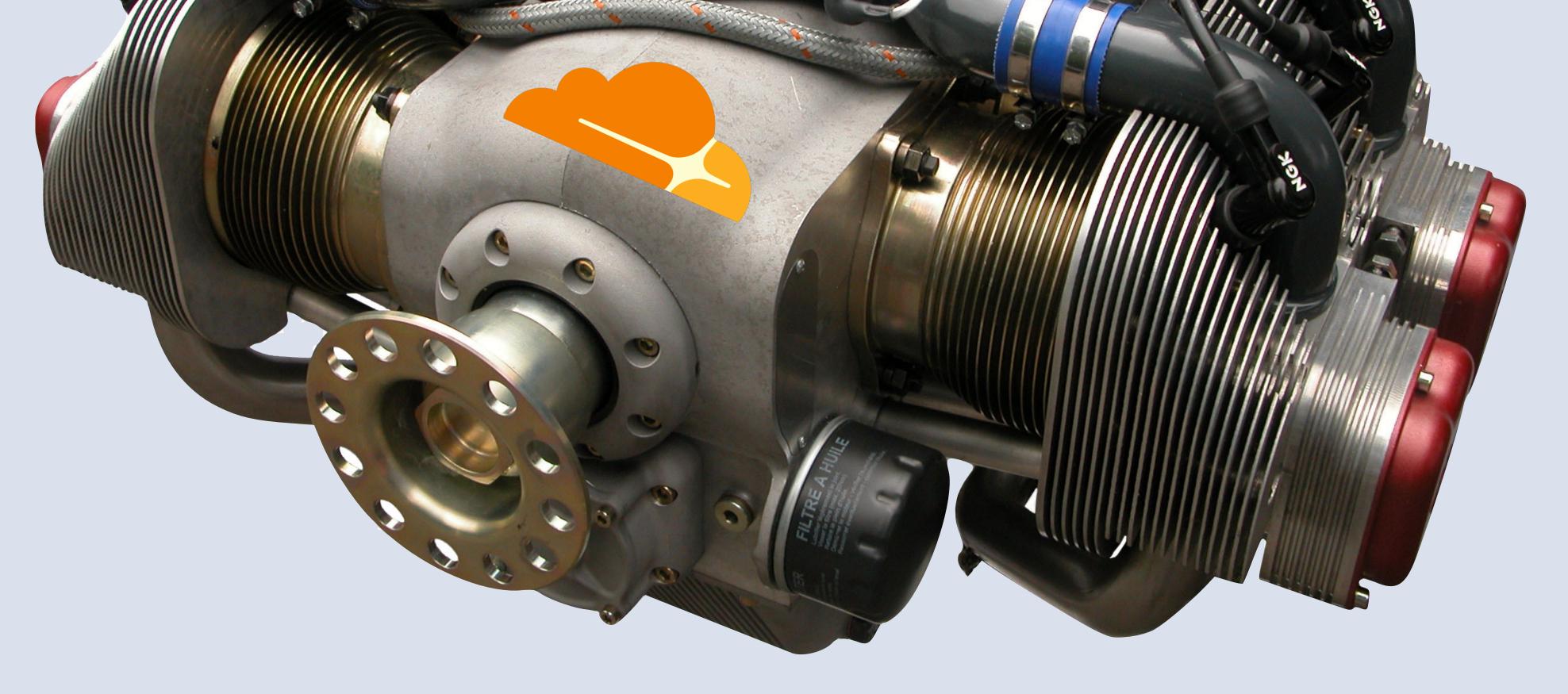
# Globally Synchronized Database



#### Constraints

# Bandwidth-limited, Broadcast





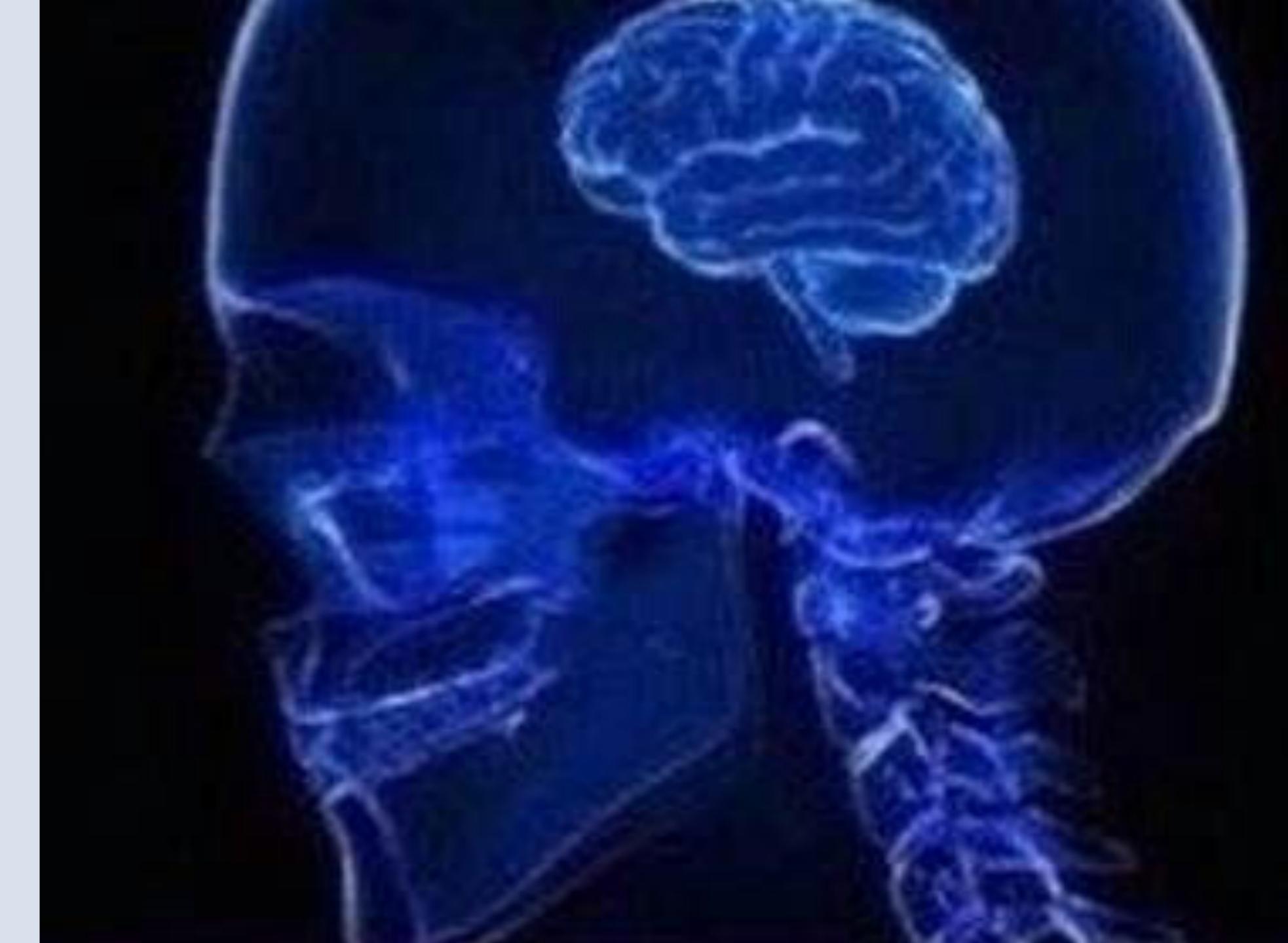
#### Identity-based provisioning system

#### **High-latency fallback**

#### **Broadcast database of keys**

#### Symmetric Cryptography





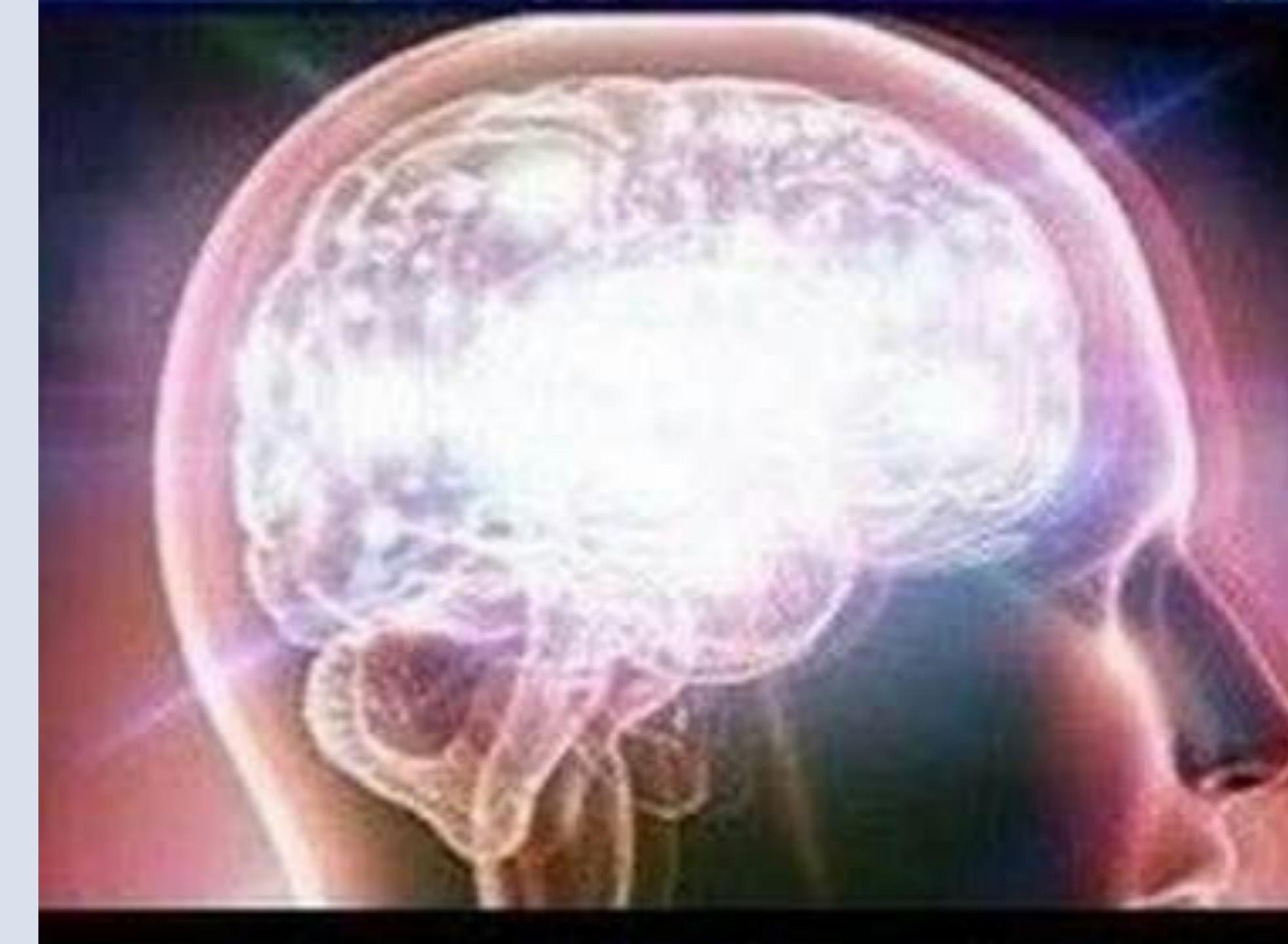
#### Asymmetric Cryptography





#### Pairing-based Cryptography





### Fully Homomorphic Encryption



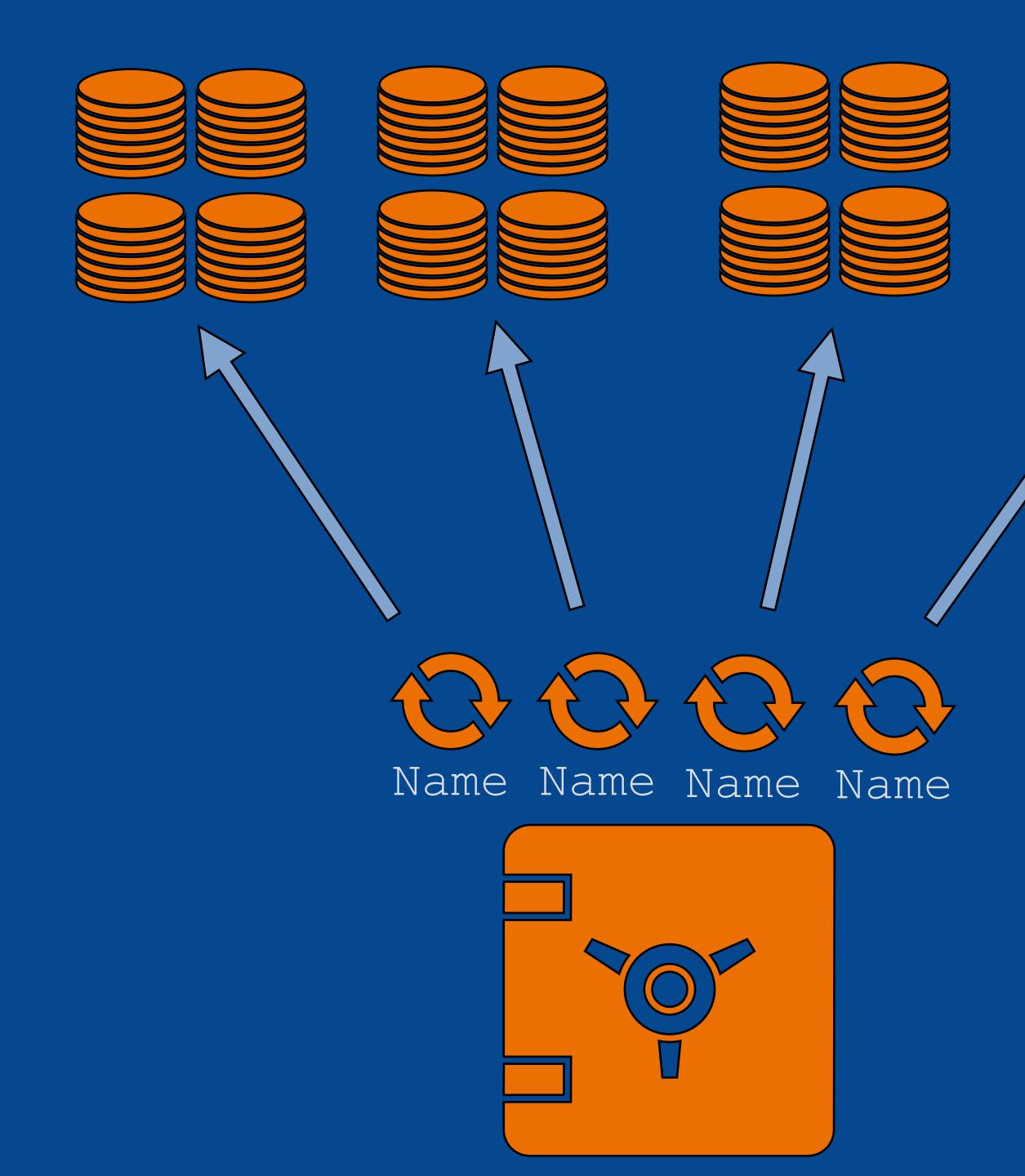


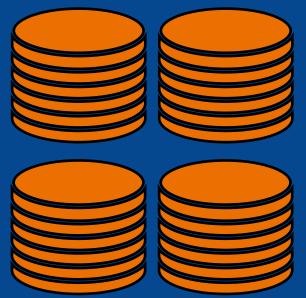
# Identity-based encryption

**Public Key:** used to encrypt data to any identity (like "machine2") Master Key: provisions private keys to identities **Private Key**: decrypts ciphertext

Allows encryption to identities even if they don't have a key yet







#### Participants

#### Private Keys

Extract

#### Master Key







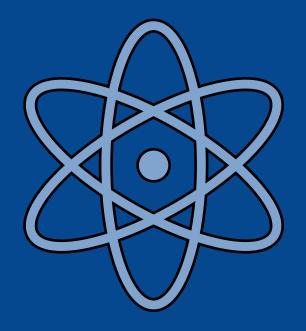


#### Encrypt

#### Ciphertext

#### JAC



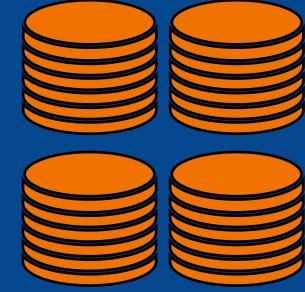












# Bilinear Pairings

### e: $G_1 \times G_2 \longrightarrow G_T$

#### First functional IBE by Boneh & Franklin (2001)



- $e(P + Q, R) = e(P, R) \cdot e(Q, R)$
- $e(P, Q + R) = e(P, Q) \cdot e(P, R)$

# Identity-based broadcast encryption

Public Key: used to encrypt data to any number of identities up to k
Master Key: provisions private keys
Private Key: decrypts ciphertext



# Identity-based revocation

Public Key: used to encrypt data to all identities except for k
Master Key: provisions private keys
Private Key: decrypts ciphertext



# IBBE and IBR with short ciphertexts Delerableé (2007) Attrapadung, Libert, de Panafier

- Master Key: constant
- Public Key: linear in k
- Private Key: constant
- Ciphertext: constant



Attrapadung, Libert, de Panafieu (2010)

- Master Key: constant
- Public Key: linear in k
- Private Key: linear in k
- Ciphertext: constant

# **Barreto-Naehrig Curves** e: $E(F_p) \times E'(F_{p2}) \longrightarrow F_{p12}$ **BN256**

128-bit security level\*

implementation in Go by Adam Langley

10x speedup by Brendan McMillion on x86\_64

faster than network round-trip from Zürich to Geneva



## Cloudflare IBBE and IBR with BN256 Identity (IBBE) Broadcast (IBR)

- Master Key: 226B
- Public Key: k64B + 578B
- Private Key: k64B + 64B
- Ciphertext: 192B (batching)



- Master Key: 64B
- Public Key: k64B + 384B
- Private Key: k64 + 192B
- Ciphertext: 192B

# Simplified Geo Key Manager

- 1. Each location is provisioned a private key with its name
- 2. Customer: "I want my TLS key in Zürich and New York"
- 3. Encrypt TLS key to the name of those locations
- 4. Distribute encrypted key + "available in Zürich or New York"
- 5. When a connection comes in
  - a. Decrypt key with location's private key, or
  - b. Connect to Zürich or New York with Keyless SSL



Whitelist

Put keys in multiple chosen locations Option to put keys in "new" locations based on region Blacklist

Put keys in region, but exempt specific location



### **Desired Semantics**

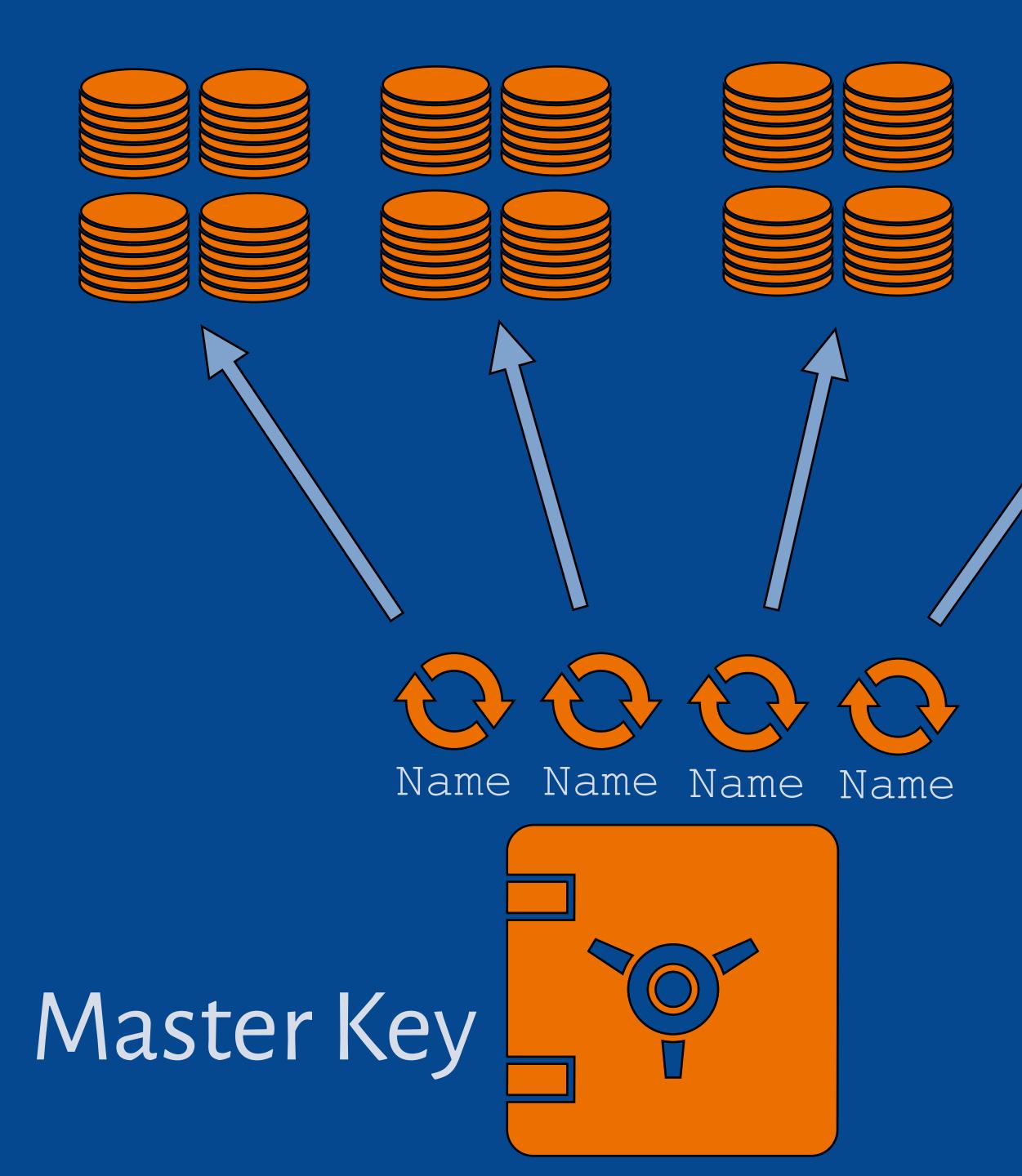
# Key Encapsulation

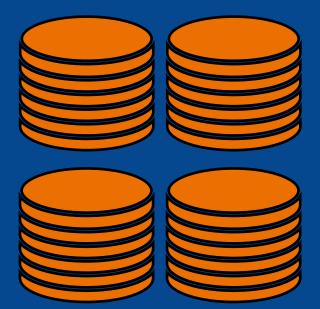
- Encrypt TLS key with a Key Encryption Key (KEK)
  - Split KEK in two
  - (e.g. KEK = KEK1 ⊕ KEK2)

KEM(kek2) for b KEM(kek2) for w



- KEM(kek1) for regions
- KEM(kek2) for blacklisted locations
- KEM(kek) for whitelisted locations





### Edge Machines

#### Private Keys

Extract

#### Provisioning Server



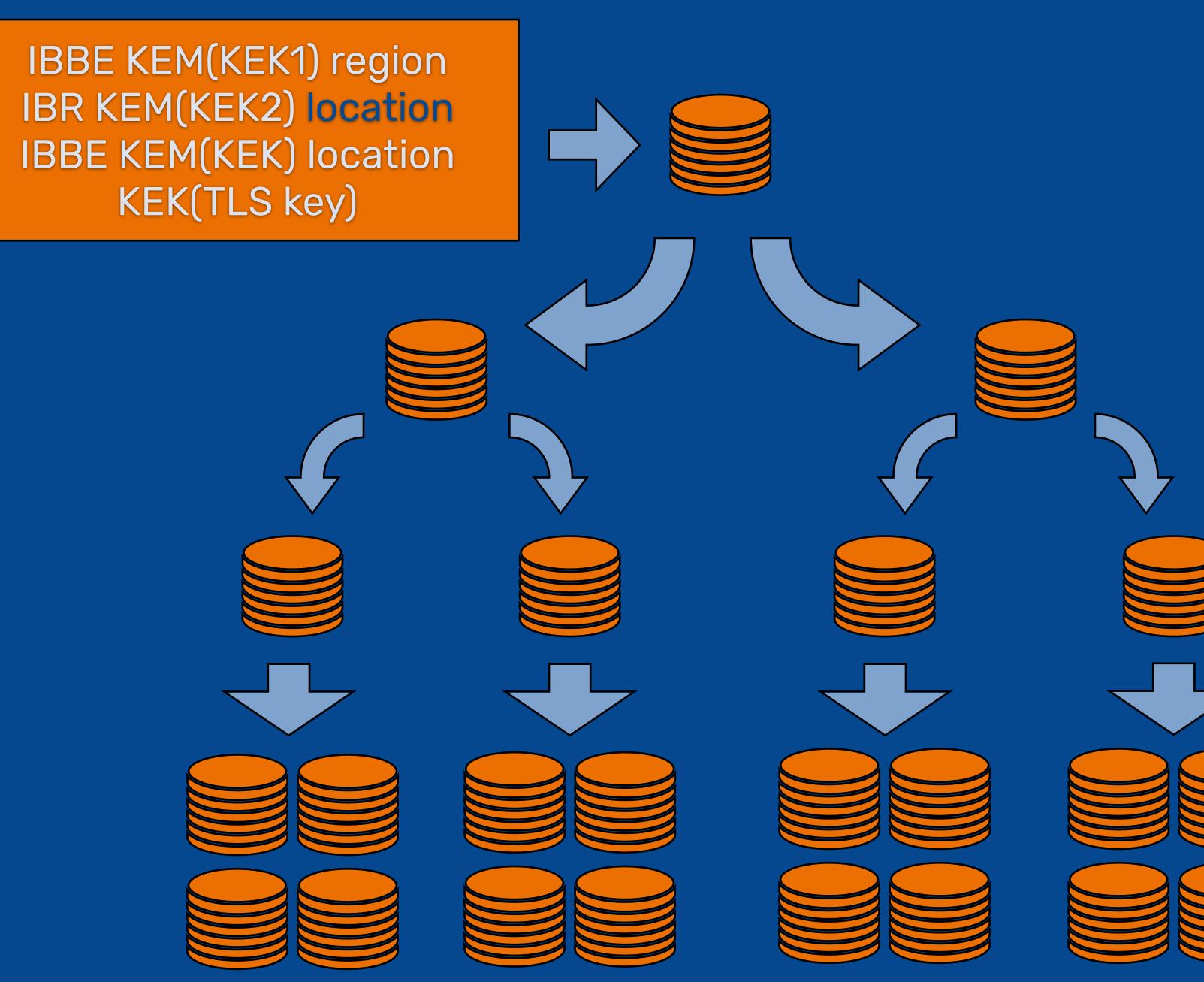




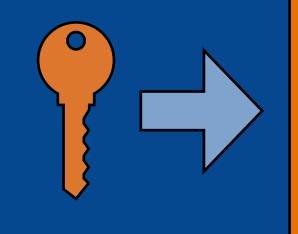


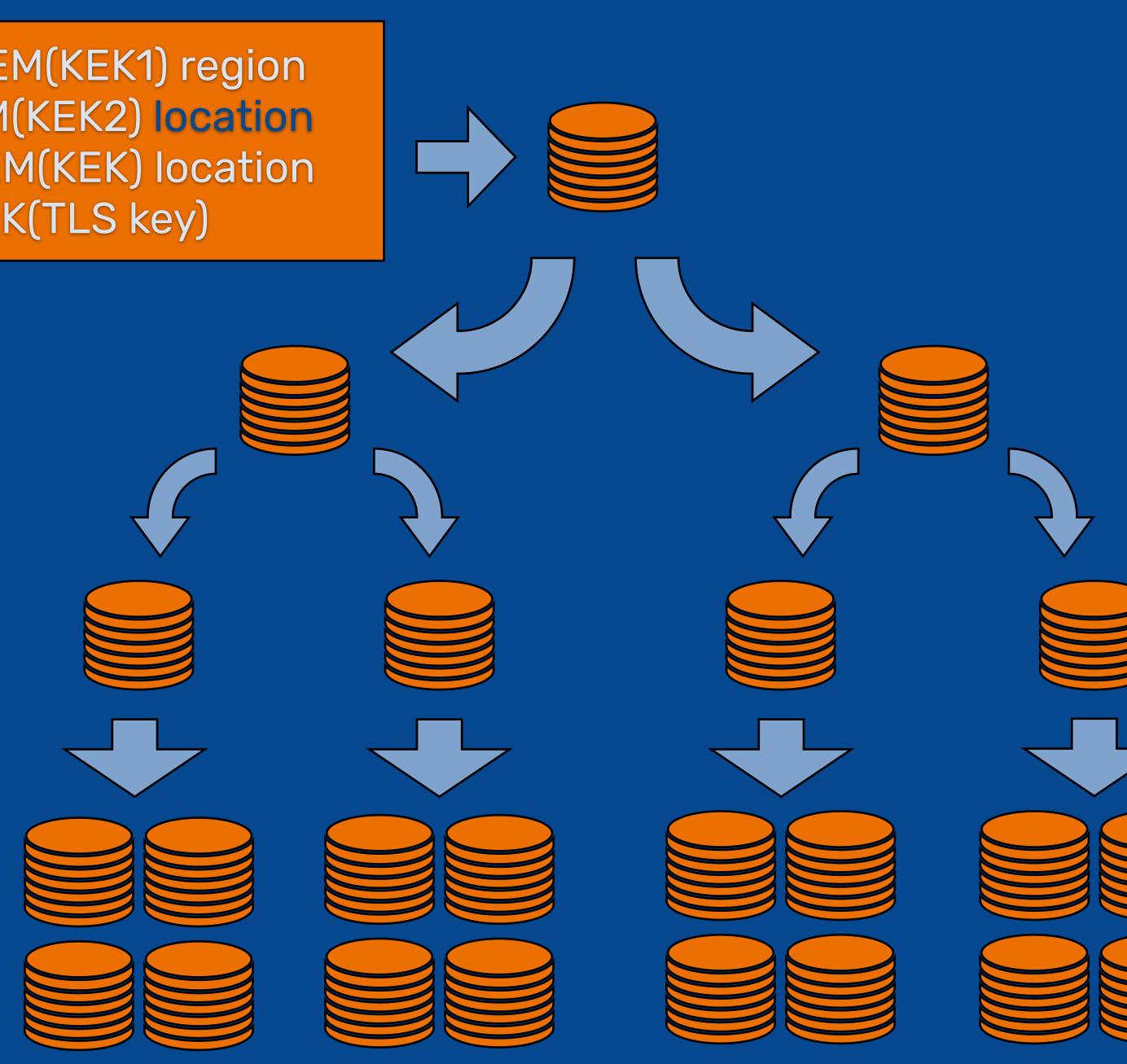




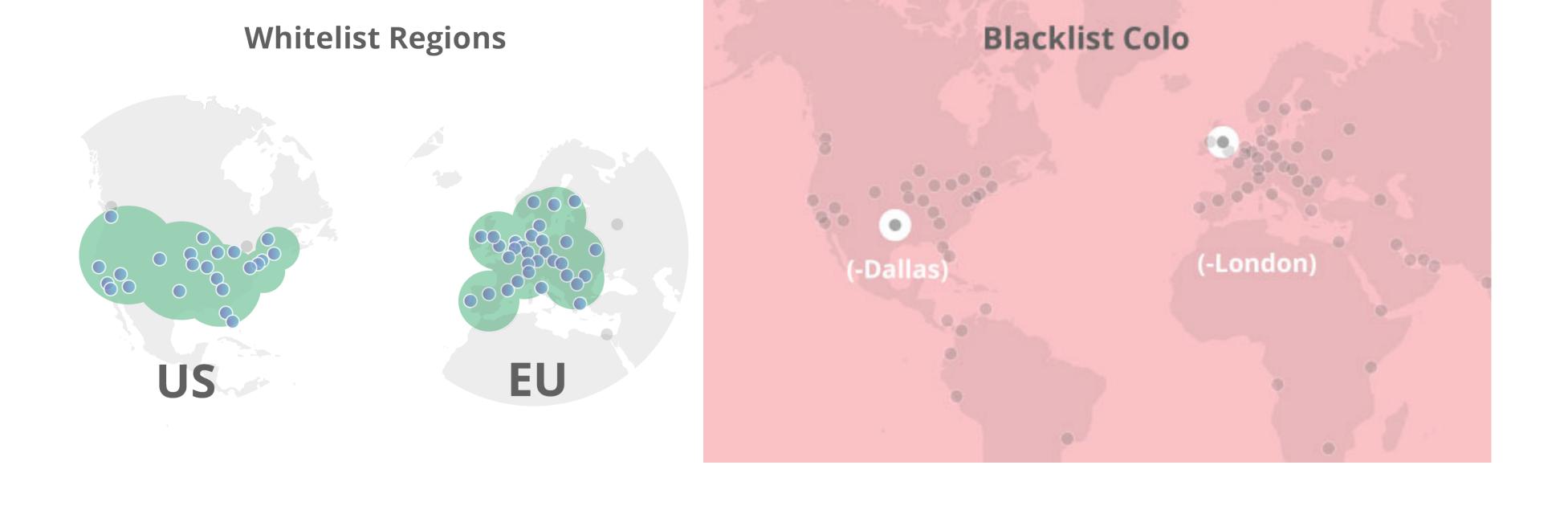


# Upload



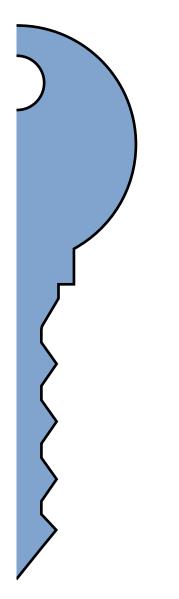


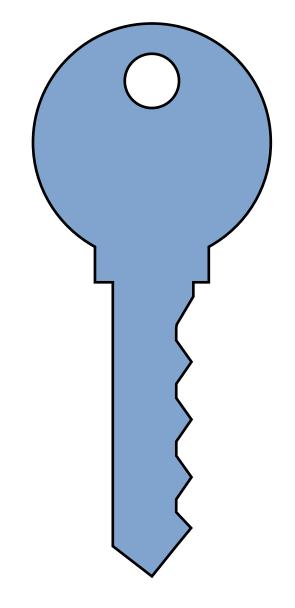


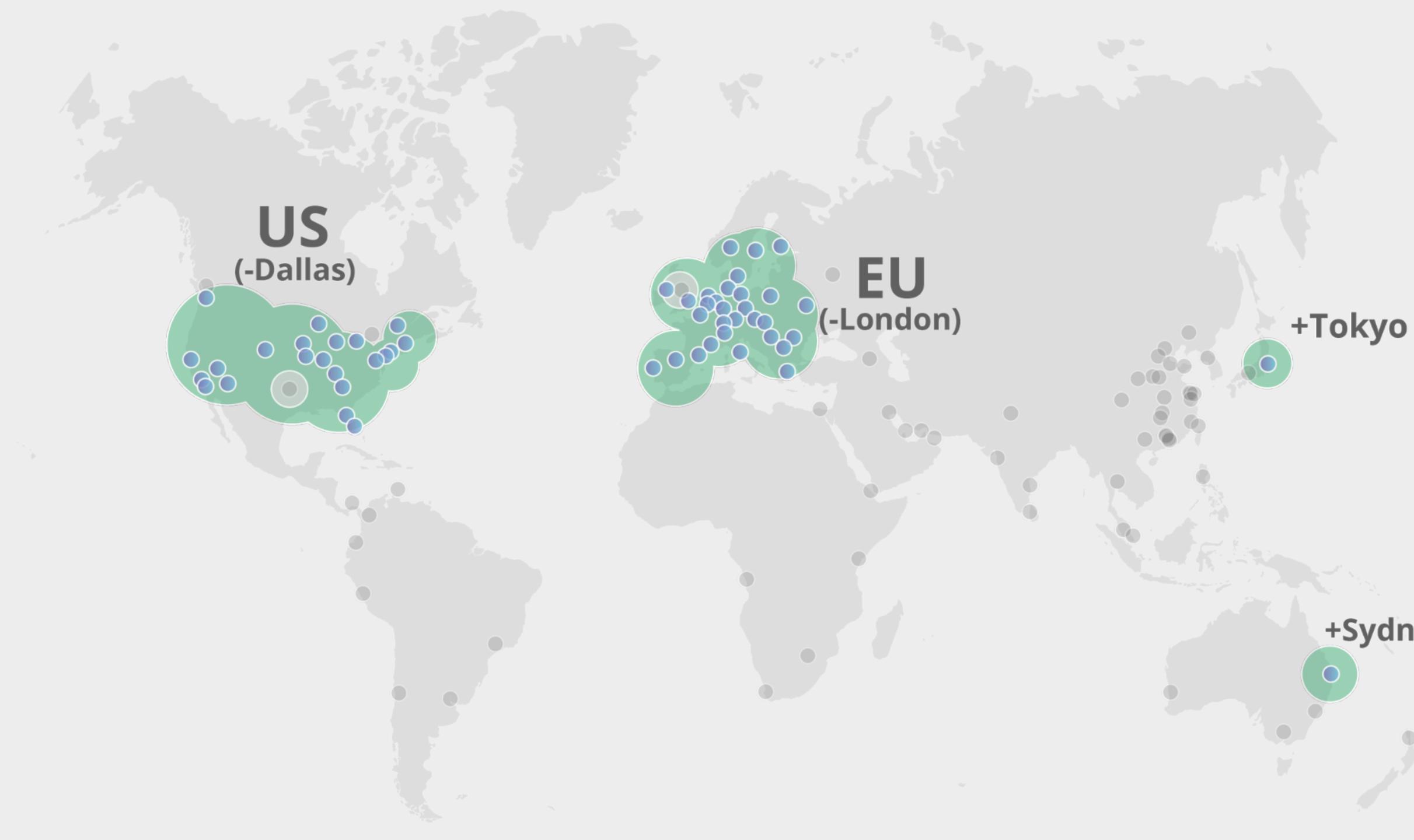


#### Whitelist Colo

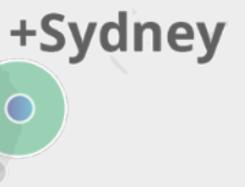














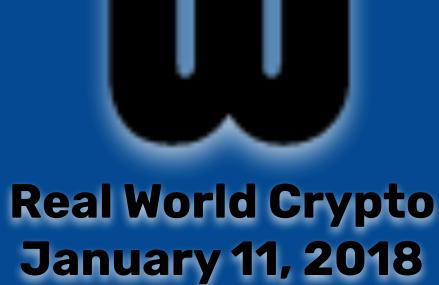
# **Geographically Distributed** Key Management

With cryptographically-enforced access control



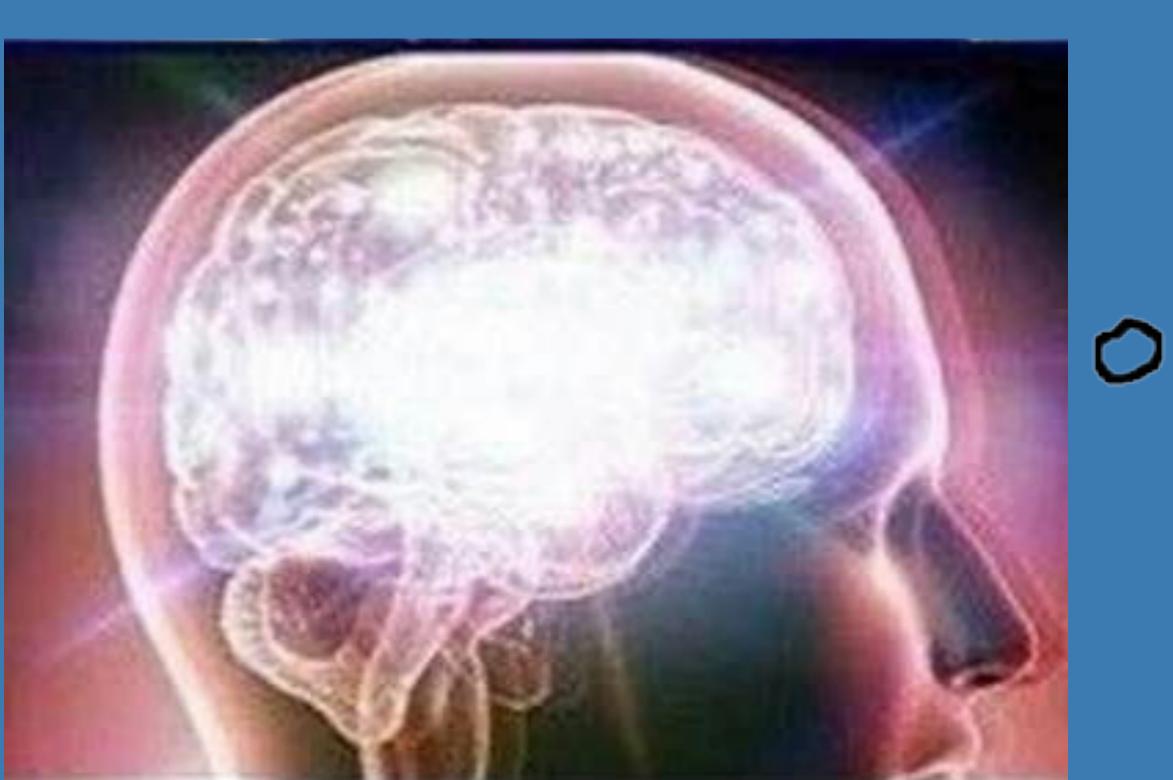


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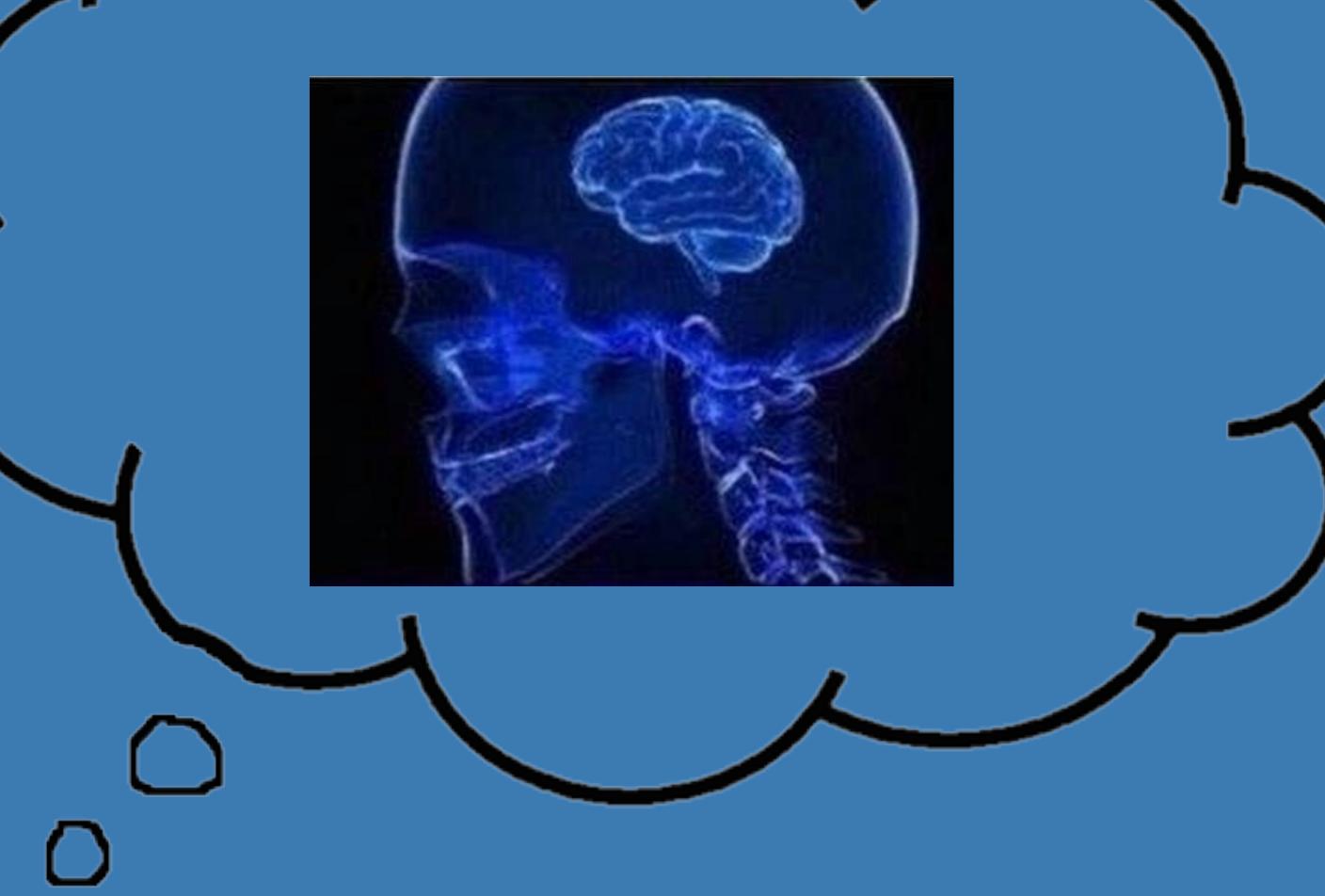


#### One pairing per symmetric key





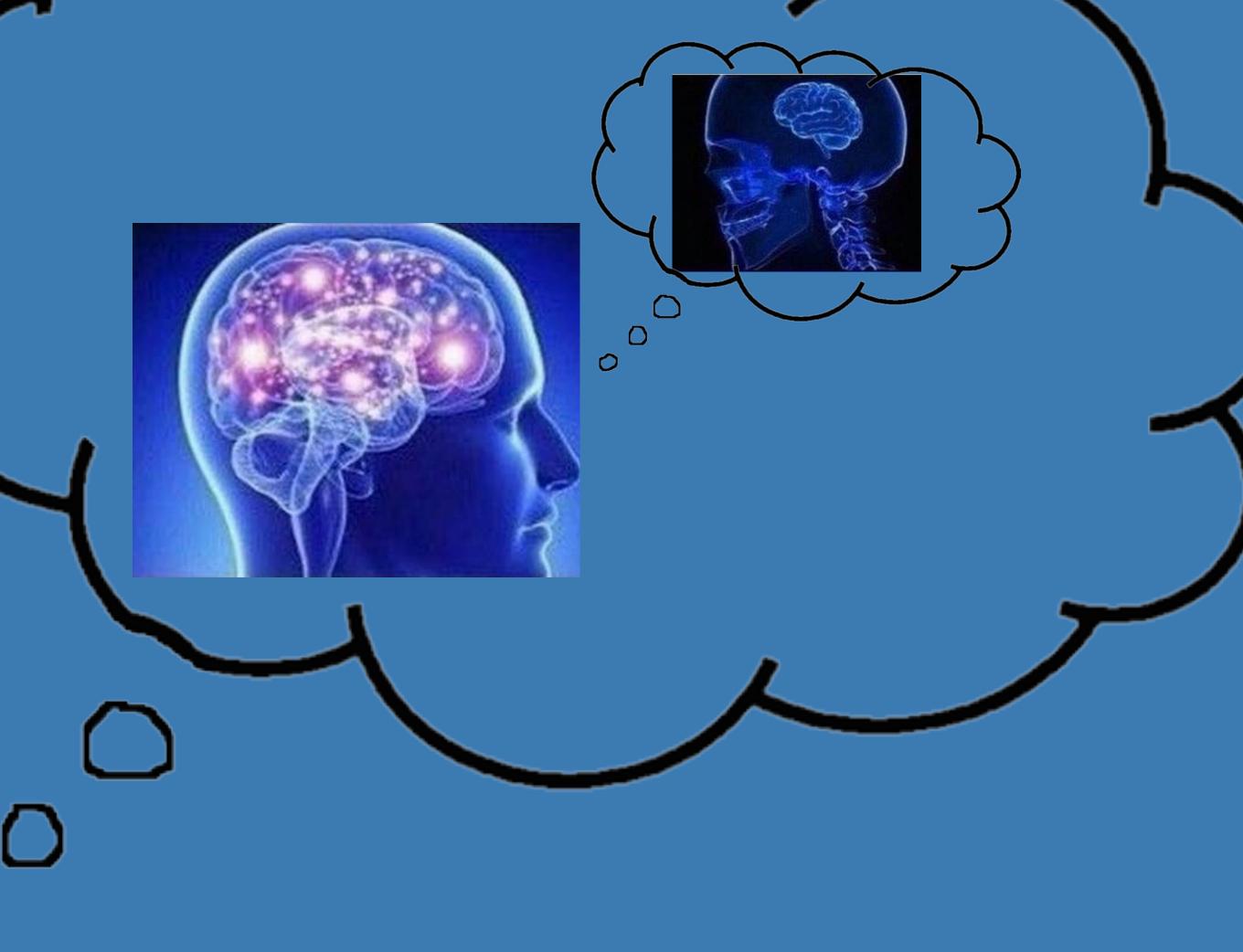




One pairing per Diffie-Hellman public key

One key exchange per symmetric key





#### For each config, generate scalars a, b, c

IBBE KEM(a) region, aP IBR KEM(b) location, bP IBBE KEM(c) location, cP

> decrypt c, compute KEK escrow c(dP), decrypt KEK or decrypt a and b and compute KEK = (a+b)dP

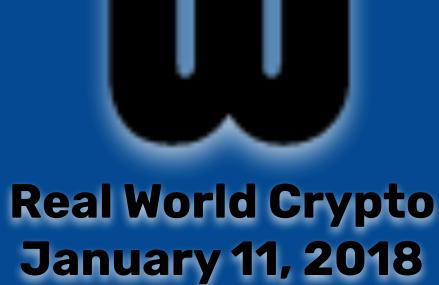
# Share KEMs between keys

For each TLS key, generate scalar d compute KEK = d(aP+bP) KEK escrow = d(cP)

> KEK(private key) KEK escrow(KEK) dP



#### **Geo Key Manager Nick Sullivan** (@grittygrease) Brendan McMillion





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